

# Much ADO about Failures: A Fault-Aware Model for Compositional Verification of Strongly Consistent Distributed Systems

**Wolf Honoré**<sup>1</sup>   **Jieung Kim**<sup>1</sup>   **Ji-Yong Shin**<sup>2</sup>   **Zhong Shao**<sup>1</sup>

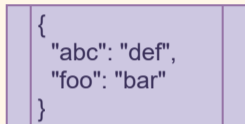
<sup>1</sup>Yale University

<sup>2</sup>Northeastern University

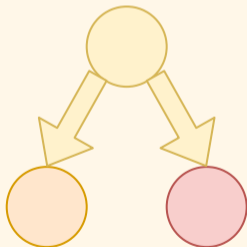
OOPSLA 2021

# Goal

**Application**  
Key-Value Store



**Implementation**  
Multi-Paxos



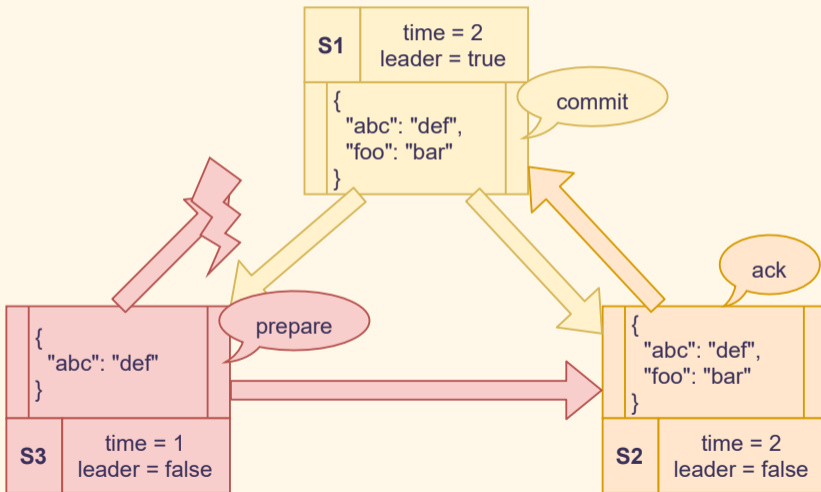
**Theorem** `KV_correct` : correct KV.

**Proof.**

...

**Qed.**

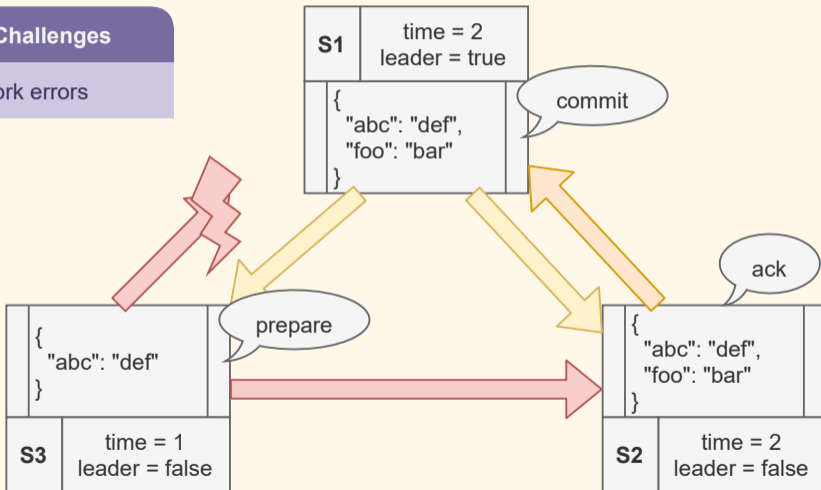
# Network-Based Models Too Complex



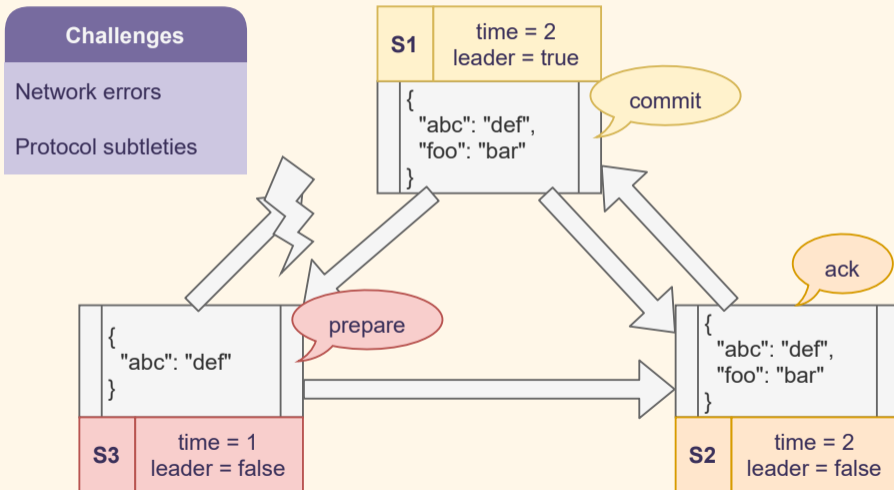
# Network-Based Models Too Complex

## Challenges

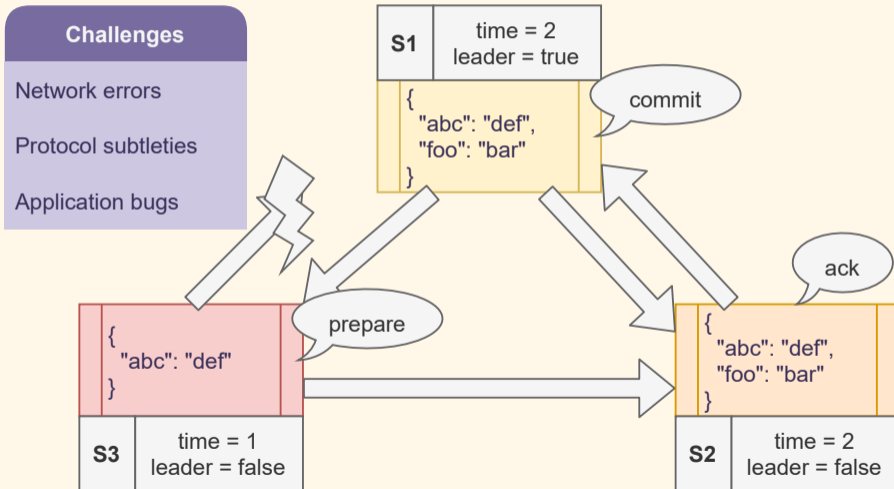
Network errors



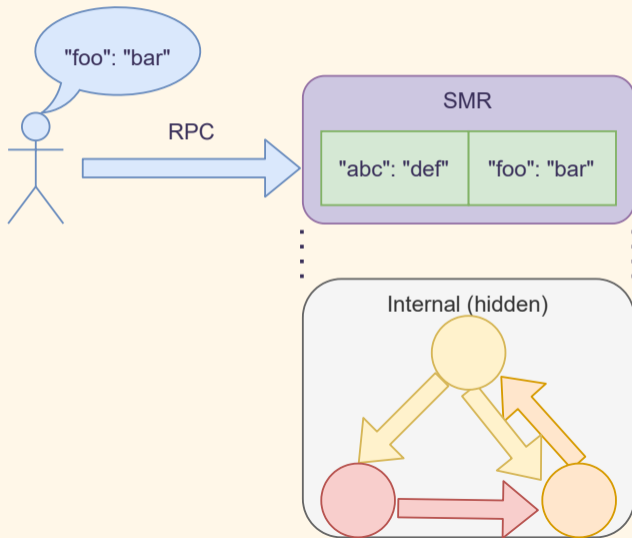
# Network-Based Models Too Complex



# Network-Based Models Too Complex

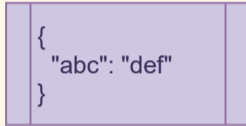


# State Machine Replication Too Abstract

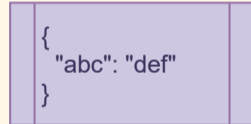


# Partial Failures

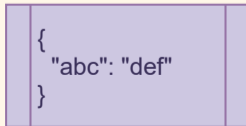
**S1**



**S2**

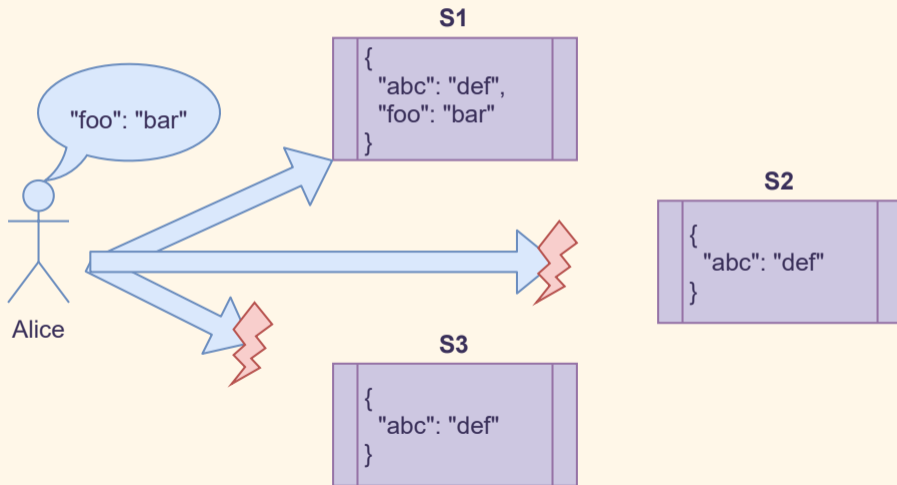


**S3**

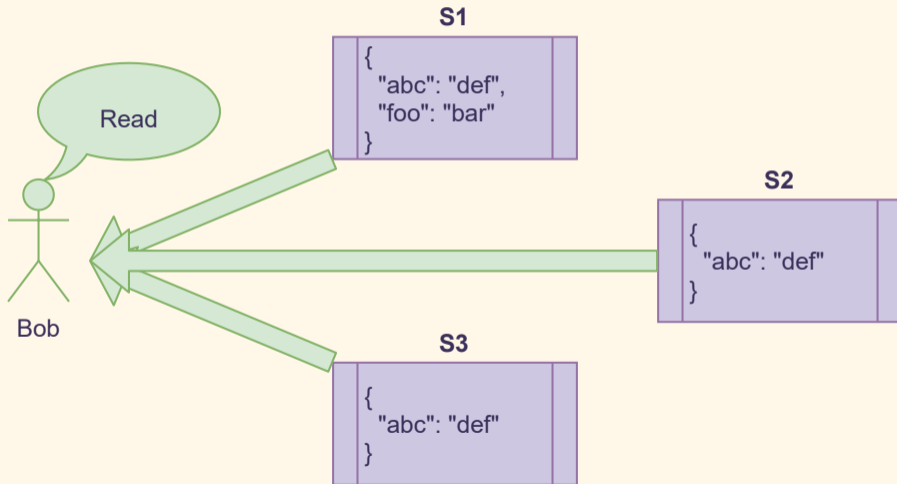




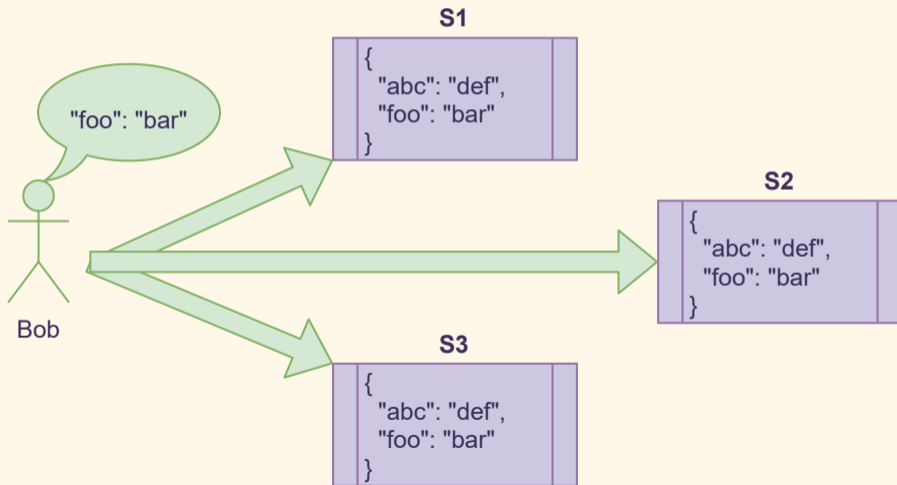
# Partial Failures



# Partial Failures



# Partial Failures



# Partial Failures are Important

*Partial failure is a central reality of distributed computing. [...] Being robust in the face of partial failure requires some expression at the interface level.*  
(Jim Waldo. *A Note on Distributed Computing*. 1994)

- ▶ Unavoidable feature unique to distributed systems.
- ▶ Influence with all aspects of distributed protocols (e.g., leader election and reconfiguration).
- ▶ Can be used for performance optimizations.
  - ▶ TAPIR (SOSP '15): Transactions with out-of-order commits.
  - ▶ Speculator (SOSP '05): Speculative distributed file system.

# A Sweet Spot?

State Machine Replication

- ✗ Hides partial failures.
- ✓ Abstracts protocol details.

?

- ✓ Shows partial failures.
- ✓ Abstracts protocol details.

Network-Based Models

- ✓ Shows partial failures.
- ✗ Blends protocol and application logic.

# Contributions

State Machine Replication

ADO Model

Network-Based Models

- ▶ ADO (atomic distributed object) model: a **fault-aware** and **compositional** abstraction.

# Contributions

State Machine Replication

ADO Model

Network-Based Models

- ▶ ADO (atomic distributed object) model: a **fault-aware** and **compositional** abstraction.
- ▶ Advert: an end-to-end Coq verification framework.
- ▶ Several verified case studies, including a lock-free key-value store, and Two-Phase Commit with replicated resource managers.

# Contributions

State Machine Replication

ADO Model

Network-Based Models

- ▶ ADO (atomic distributed object) model: a **fault-aware** and **compositional** abstraction.
- ▶ Advert: an end-to-end Coq verification framework.
- ▶ Several verified case studies, including a lock-free key-value store, and Two-Phase Commit with replicated resource managers.
- ▶ Refinement with several Paxos variants, Chain Replication.
- ▶ Refinement with multi-Paxos C implementation.



# Contributions

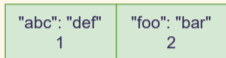
State Machine Replication

ADO Model

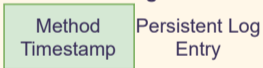
Network-Based Models

- ▶ ADO (atomic distributed object) model: a **fault-aware** and **compositional** abstraction.
- ▶ Advert: an end-to-end Coq verification framework.
- ▶ Several verified case studies, including a lock-free key-value store, and Two-Phase Commit with replicated resource managers.
- ▶ Refinement with several Paxos variants, Chain Replication.
- ▶ Refinement with multi-Paxos C implementation.

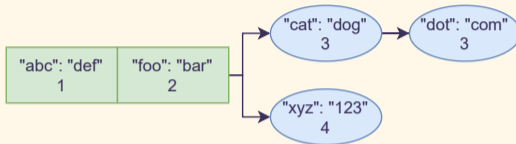
# ADO State



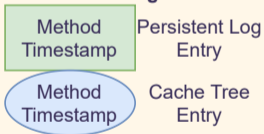
## ADO Legend



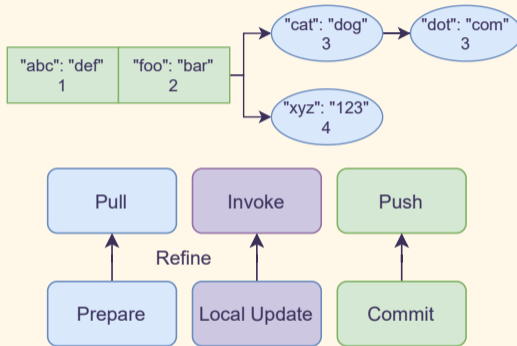
# ADO State



## ADO Legend

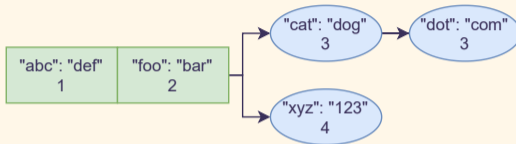


# ADO Operations

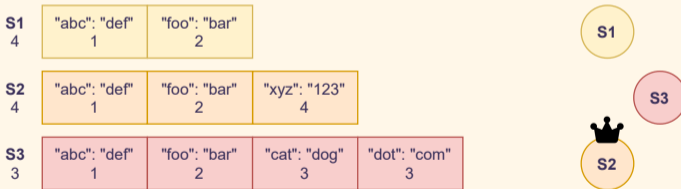


# ADO Operations

ADO

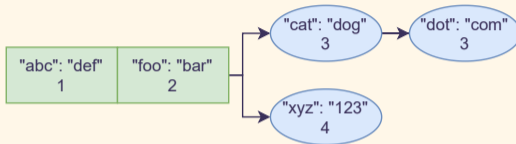


Multi-Paxos

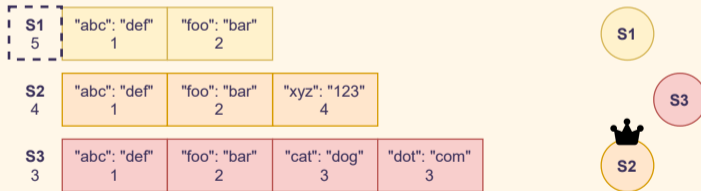


# Pull

ADO

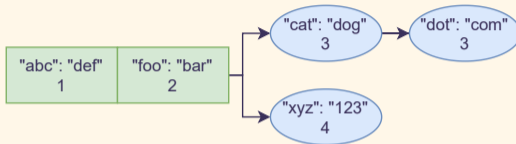


Multi-Paxos

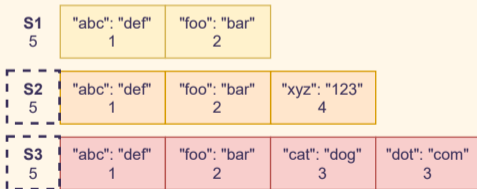


# Pull

ADO

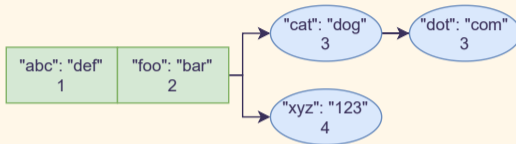


Multi-Paxos

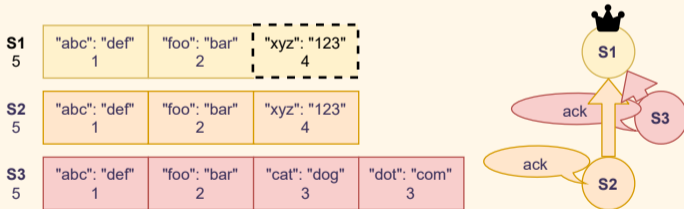


# Pull

ADO

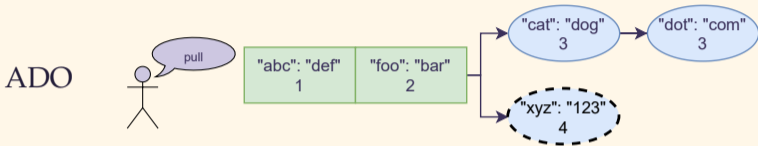


Multi-Paxos

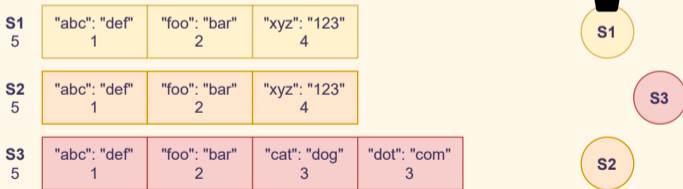




# Pull



## Multi-Paxos

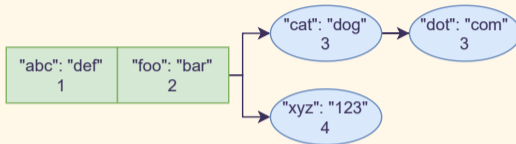


### Pull

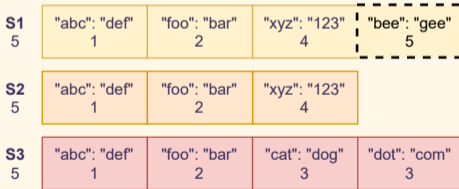
Get permission to update and select a starting point in the cache tree.

# Invoke

ADO

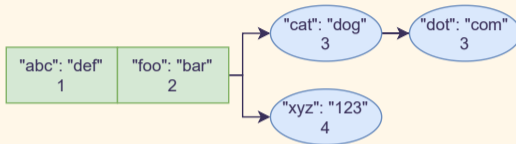


Multi-Paxos



# Invoke

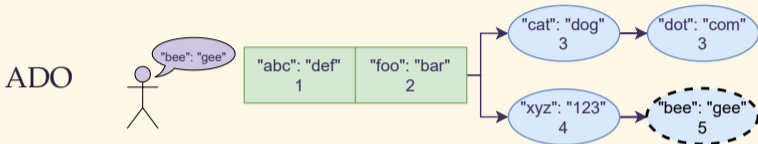
ADO



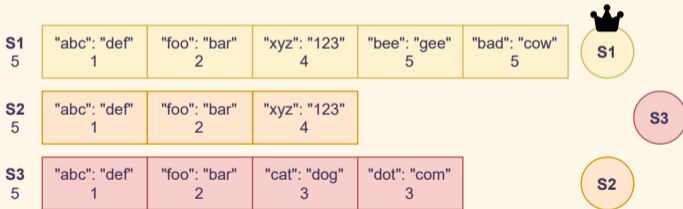
Multi-Paxos



# Invoke



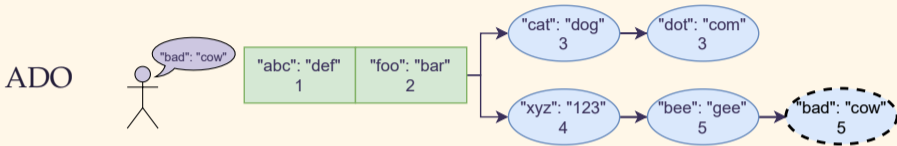
## Multi-Paxos



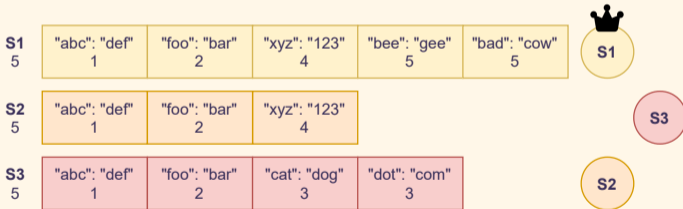
### Invoking a Method

Add a new entry to the cache tree.

# Invoke



## Multi-Paxos

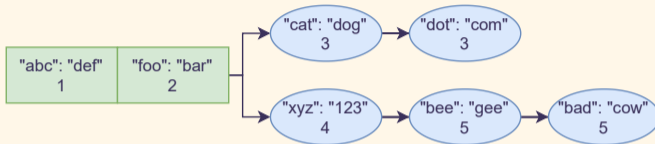


## Invoking a Method

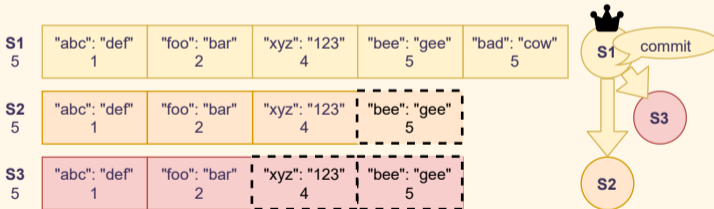
Add a new entry to the cache tree.

# Push

ADO

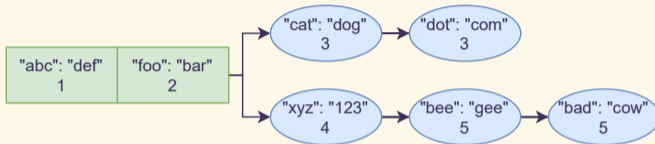


Multi-Paxos



# Push

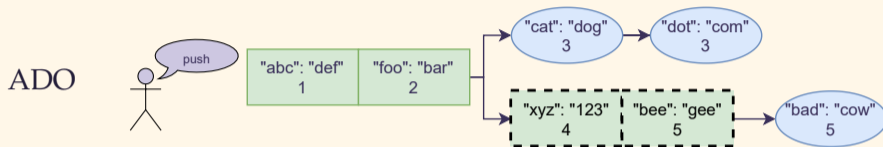
ADO



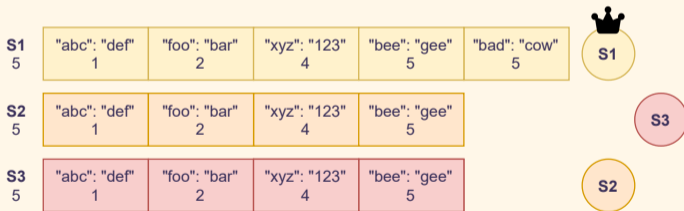
Multi-Paxos



# Push



## Multi-Paxos

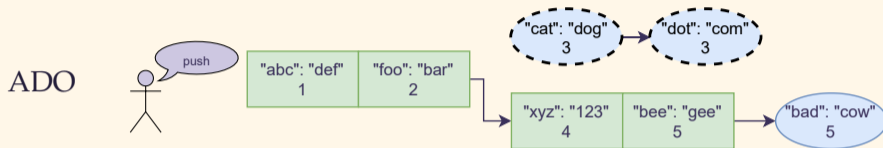


## Push

Move committed methods into the log and prune stale states from the tree.



# Push



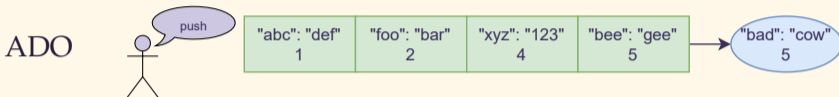
## Multi-Paxos



## Push

Move committed methods into the log and prune stale states from the tree.

# Push



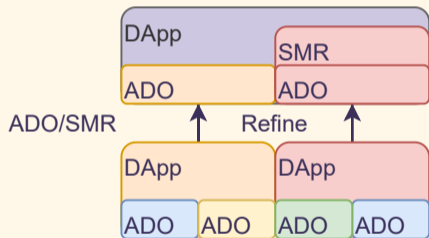
## Multi-Paxos



### Push

Move committed methods into the log and prune stale states from the tree.

# Distributed Applications



# Distributed Applications

```
ADO KV {  
  shared kv : [string * int] := [];  
  method set(k, v) { this.kv[hash(k)] := (v, len(v)); }  
  method get(k) { return this.kv[hash(k)][0]; }  
  method getmeta(k) { return this.kv[hash(k)][1]; }  
}
```

# Distributed Applications

```
ADO DVec[T] {  
  shared data : [T] := [];  
  method insert(idx, x) { this.data[idx] := x; }  
  method get(idx) { return this.data[idx]; }  
}
```

```
ADO DLock {  
  shared owner : option N := None;  
  method tryAcquire() { ... }  
  method release() { ... }  
}
```

```
DApp KVLock(lk: DLock, data: DVec[string], meta: DVec[int]) {  
  proc set(k, v) {  
    ... /* acquire, set data, set meta, release */  
  }  
  ... /* get, getmeta */  
}
```

# Distributed Applications

```
DApp KVLock(lk: DLock, data: DVec[string], meta: DVec[int]) {  
  proc set(k, v) {  
    lk.pull();  
  
  }  
}
```

# Distributed Applications

```
DApp KVLock(lk: DLock, data: DVec[string], meta: DVec[int]) {  
  proc set(k, v) {  
    while (lk.pull() == FAIL) {}  
  
  }  
}
```

# Distributed Applications

```
DApp KVLock(lk: DLock, data: DVec[string], meta: DVec[int]) {  
  proc set(k, v) {  
    while (lk.pull() == FAIL) {}  
    ok := lk.tryAcquire();  
  
  }  
}
```



# Distributed Applications

```
DApp KVLock(lk: DLock, data: DVec[string], meta: DVec[int]) {  
  proc set(k, v) {  
    while (lk.pull() == FAIL) {}  
    ok := lk.tryAcquire();  
    while (lk.push() == FAIL) {}  
    if (!ok) { return; }  
    /* ... */  
  }  
}
```

# Method Calling Semantics

```
DApp KVLockAbort(lk: DLock, data: DVec[string], meta: DVec[int]) {  
  proc set(k, v) {  
    if (lk.pull() == FAIL) { return; }  
    ok := lk.tryAcquire();  
    if (lk.push() == FAIL) { return; }  
    if (!ok) { return; }  
    /* ... */  
  }  
}
```

# Method Calling Semantics

```
DApp KVLockRetry(lk: DLock, data: DVec[string], meta: DVec[int]) {  
  proc set(k, v) {  
    for retry in 0..N {  
      if (lk.pull() == FAIL) { continue; }  
      ok := lk.tryAcquire();  
      if (lk.push() == FAIL) { continue; }  
      if (!ok) { continue; }  
    }  
    if (retry == N) { return; }  
    /* ... */  
  }  
}
```

# Method Calling Semantics

```
obj.m()! :=  
  while (obj.pull() == FAIL) {}  
  obj.m();  
  while (obj.push() == FAIL) {}
```

```
DApp KVLock(lk: DLock, data: DVec[string], meta: DVec[int]) {  
  proc set(k, v) {  
    ok := lk.tryAcquire()!;  
    if (!ok) { return; }  
    data.insert(hash(k), v)!;  
    meta.insert(hash(k), len(v))!;  
    lk.release()!;  
  }  
}
```

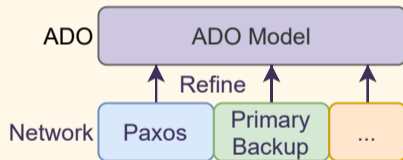
## Non-Standard Method Calls

```
DApp TM(rm_1: RM, ..., rm_n: RM) {  
  proc init() { // Must be called once when TM starts  
    for rm in [this.rm_1, ..., this.rm_n] {  
      while (rm.pull() == FAIL) {} // pull once up front  
    }  
  }  
  proc collect_decisions(tx) {  
    for rm in [this.rm_1, ..., this.rm_n] {  
      rm.prepare(tx); // No pull needed  
      for i in 0..MAX_TRY {  
        res := rm.push(); // Only try up to MAX_TRY  
        if (res != FAIL) { break; }  
      }  
      // Short-circuit on failure  
      if (res == NO || res == FAIL) { tx.decision := ABORT; break; }  
    }  
  }  
}
```

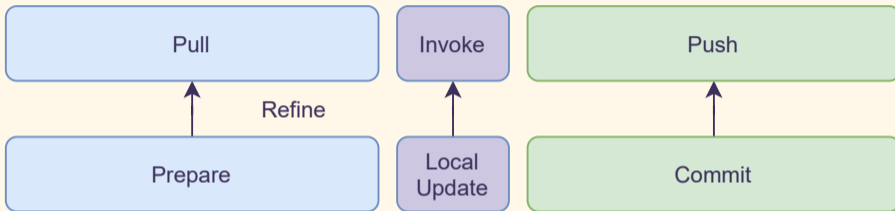
## Non-Standard Method Calls

```
DApp TM(rm_1: RM, ..., rm_n: RM) {  
  proc init() { // Must be called once when TM starts  
    for rm in [this.rm_1, ..., this.rm_n] {  
      while (rm.pull() == FAIL) {} // pull once up front  
    }  
  }  
  proc collect_decisions(tx) {  
    for rm in [this.rm_1, ..., this.rm_n] {  
      rm.prepare(tx); // No pull needed  
      for i in 0..MAX_TRY {  
        res := rm.push(); // Only try up to MAX_TRY  
        if (res != FAIL) { break; }  
      }  
      // Short-circuit on failure  
      if (res == NO || res == FAIL) { tx.decision := ABORT; break; }  
    }  
  }  
}
```

# Connection with Distributed Protocols

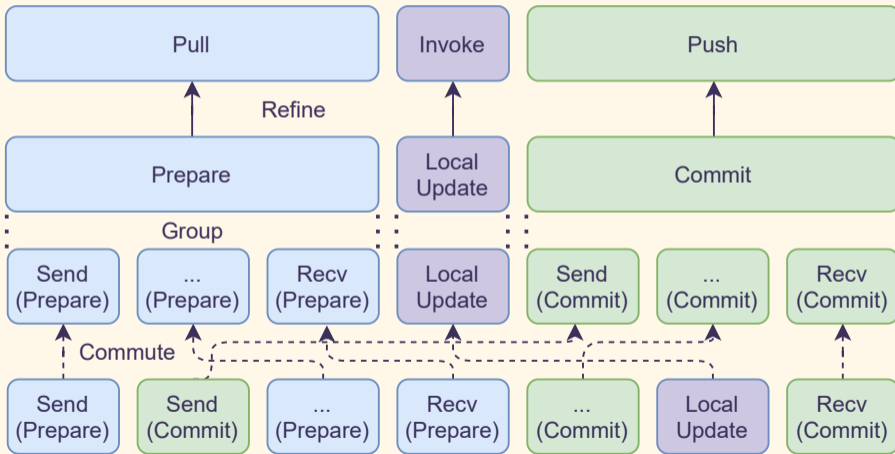


# Refinement

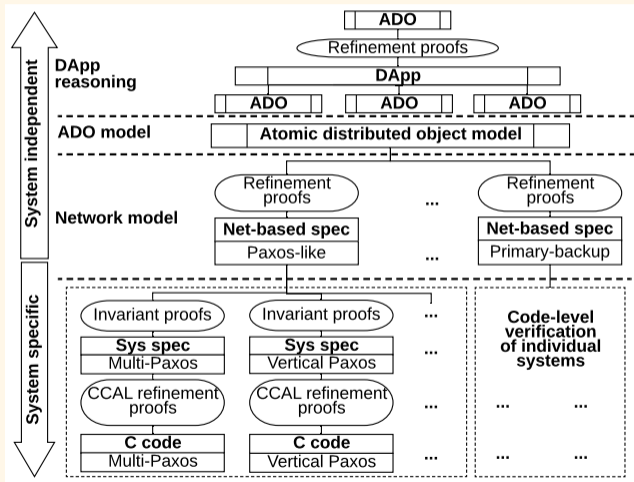




# Refinement



# Specification and Proof Effort



Component	LOC
KVLock	646
KVLockFree	359
2PC	559
Paxos-like	5K
Chain Replication	2K
Shared Libraries	11K
Single-Paxos	77
Multi-Paxos	87
Vertical Paxos	97
CASPaxos	78

Code available at <https://zenodo.org/record/5476274>.

# Conclusion

- ▶ ADO Model: A novel, fault-aware, compositional distributed system abstraction.
- ▶ Advert: Coq framework for single- and multi-ADO reasoning.
- ▶ End-to-end guarantees with refinement.
- ▶ High-level behavior independent of underlying protocol.